PL Evaluation and Typing Systems

Principles of Programming Languages

Colorado School of Mines

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Evaluating a Programming Language

Evaluation Metrics

Evaluating programming languages based on:

Writability:	*How easy is it to write good code?*
Readability:	*How easy is it to read well written code? Is the language easy enough to learn?*
Reliability:	*What features does the language provide to make sure our code works as it is supposed to?*
Feasibility:	*Does an interpreter or compiler actually exist for the platform we need to use? Is it fast enough for our application?*

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A System of Trade-Offs

Often times, adding features which improve one metric can harm another metric. Examples to come...

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Simplicity

The overall simplicity of a language plays a large role in both writability and readability.

For example, these features are *non-simple*:

- **Feature Multiplicity:** 👍 Writability, 👎 Readability
- **Operator Overloading:** 👍 Writability, 👎 Readability
- **Large Grammars:** 👍 Writability, 👎 Readability

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Simplicity can be carried too far

Assembly languages and esoteric languages generally aren't considered very writable or readable.

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Orthogonality: how consistent is the language with itself?

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Example of a lack of orthogonality (C/C++)

Arrays can contain data of any type, including pointers.
Unless it's a function pointer.
But you can wrap that function pointer in a `struct` and you should be fine.

Impacts of poor orthogonality: poor readability, poor writability, and potentially reduced reliability.

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Abstraction

Abstraction: *The ability to define and use complicated structures and operations in a way that allows implementation to be ignored.*

Examples:

- **Functions:** Simplest form of abstraction. Often taken for granted, but gives us easy recursion.
- **Heap Memory:** Imagine trying to create a large unbalanced binary tree in a single-dimensional array.
- **Generics:** Allows us to define operations that apply to multiple data types without reimplementing for each type.

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With your learning group...

What other kinds of PL-level abstractions can you name?

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What other kinds of PL-level abstractions can you name?

Good Abstractions: 👍 Writability, 👍 Readability, 👍 Reliability

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With your learning group...

1. Can abstractions be taken *too* far? If so, give an example?
2. Where else do we find abstractions in Computer Science?

Reliability Features

Some languages come with features *designed for reliablitiy*:

- **Type Checking:** Making sure the type of data can be used with the function or operation you are calling. Independent of static/dynamic: more on this later.
- **Exception Handling:** The ability of a running program to intercept run-time errors and take corrective measures.
- **Taint Protection:** Protects the security of an application by not allowing privileged operations to be preformed on tainted data (e.g., user input from a web application).

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Some features can *harm* a language's reliability:

- **Goto:** the ability to jump to different locations in the code without restriction.
- **Aliasing:** allows two different symbolic names (variables, function names, etc.) to refer to the same data. Think pointers in C/C++.

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Expressivity: How easy is it for the programmer to express their solution to a problem in the language?

With your learning group:

1. Think of a scenario and two programming languages, where expressing a solution to the problem might be easier in one language than another.
2. If one language is less expressive than another, how might it be less **writable**?
3. If one language is less expressive than another, how might it be less **readable**?
4. If one language is less expressive than another, how might it be less **reliable**?

Be prepared to share your answers with the class.

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Type Systems

When we refer to "type systems", we aren't just talking about how you have to write the type of a variable in C, whereas you don't in Python. There's a lot that goes in:

- Static/Dynamic Typing
- Untyped Systems
- Implicit/Explicit Type Specification
- Strong/Weak Typing
- Gradual Typing
- Duck Typing -- covered later in the course

Explicit/Implicit

In a language with **explicit type specification**, the type of a variable *must* be specified:

```
int x = 10;
```

In a language with **implicit type specification**, the type of a variable need not be specified:

```
var x = 10
```

Note

Explicit/implicit has nothing to do with static/dynamic. We will talk about that in a second...

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Bindings

A **binding** refers to the association between:

- a variable and its type,
- a function and its definition,
- a type and its representation (e.g., `int` is 32-bits),
- or an operation and its symbol (e.g., multiplication is usually *)

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Binding time refers to the time at which a binding takes place.

Common Binding Times

Design time, implementation time, compile time, link time, run time

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Static Typing

In a **statically typed** language, the language design makes it possible to bind the type of any piece of data **before run time**.

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Advantages:

- No need to do type checking at run time, this can be done at compile time.
- 👍 Reliability

Disadvantages:

- Generics are needed to create operations and functions that apply to multiple types
- 👎 Writability

Dynamic Typing

In a **dynamically typed** language, the language design makes it possible to bind the type of any piece of data **during run time**. Commonly, the type of data is associated with the data itself.

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In a **dynamically typed** language, the language design makes it possible to bind the type of any piece of data **during run time**. Commonly, the type of data is associated with the data itself.

Advantages:

- Collections can be of mixed type without generics, functions can take multiple types without generics
- Types can be dynamically created at run time
- 👍 Writability

Disadvantages:

- Type checking must be done at run time; makes things slow
- 👎 Reliability

Gradual Typing

Gradual Typing can be used to refer to a language which allows optional explicit typing in a dynamically typed language.

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Untyped Systems

In an **untyped** language, data cannot be bound to a type.

Commonly, the functions and operations called on the variables determine the type.

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In an **untyped** language, data cannot be bound to a type.

Commonly, the functions and operations called on the variables determine the type.

Advantages:

- No need to do type checking, ever.
- 👍 Feasibility

Disadvanges:

- 👎 👎 👎 Reliability

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- **Type safety** means a language will not allow bits to be intepreted as the incorrect data type. For example: treating the bits of an integer as a floating point number.
- **Implicit type conversions** are when a language will automatically convert data types to allow an expression to be computed.
- **Strongly typed** programming languages are both type safe *and* do not allow implicit type conversions.
- **Weakly typed** programming languages are either not type safe *or* allow implicit type conversions.

Note

By definition, untyped languages are also weakly typed.

Given the code snippet from a fake language below:

```
int a = 10
a += 5
print(a)
```

- Explicit or implicit?
- Is it possible that the language is statically typed?
- Is it possible that the language is dynamically typed?
- Weak or strong?

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- Is it possible that the language is dynamically typed?
- Weak or strong?

Given the code snippet from a fake language below:

```
a = eval(user_input())
a += 5
print(a)
```

- Explicit or implicit?
- Is it possible that the language is statically typed?
- Is it possible that the language is dynamically typed?
- Weak or strong?

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Given the code snippet from a fake language below:

```
++++++++[>++++[>++>+++>+++>+<<<<-]>+
>+>->>+[<]<-]>>.>---.+++++++..+++.>>
.<-.<.+++.------.-----+----.>>+.>++.
```

- Explicit or implicit?
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- Is it possible that the language is dynamically typed?
- Weak or strong?

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Type Systems: Language Examples

	Strong	**Weak**
Static	Haskell, Rust, Go	C, C++
Dynamic	Python, Ruby, Java	PHP, JavaScript

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